

# Finite Summation of Integer Powers $x^p$ , Part 3

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## Abstract

We find a *direct* closed-form solution, i.e. one that does *not* require iteration, for the general case of the finite-summation-of-integer-powers problem  $S_p(N) = \sum_{k=1}^N k^p$ . Having established in Part 2 [Ebr10b] that the closed-form solution is a polynomial, the summation is here rewritten as the sum of the  $p+1$  independent monomials  $a_j N^j$  ( $1 \leq j \leq p+1$ ), where the  $a_j$  are unknown coefficients. Using the recurrence relation  $S_p(N+1) = S_p(N) + (N+1)^p$ , we obtain a linear combination of the monomials, which reduces to an easily solvable  $(p+1)$ -by- $(p+1)$  triangular linear system in the unknown coefficients  $a_j$  of the closed-form polynomial solution. Maxima and Octave/Matlab codes for directly computing the closed-form solutions are included in the Appendices.

## Finding a Closed-Form solution for $S_p(N) = \sum_{k=1}^N k^p$ using a Direct Method

### Motivation

Our goal is to obtain, without using iteration, a closed form solution for the general case of the finite-summation-of-integer-powers problem:

$$S_p(N) = \sum_{k=1}^N k^p, \text{ (where } p \in \mathbb{N} = \{0, 1, 2, \dots\}; N \geq 0) \quad (1)$$

Recall that in Part 2 [Ebr10b] of this paper, we used a method of recurrence relations to obtain the following results about  $S_p(N)$ :

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1.  $S_p(N)$  may be expressed as a linear recurrence relation that uses the  $p-1$  lower order formulas  $S_j(N)$  and  $S_j(N-1)$  ( $j = 1, \dots, p-1$ ) as follows:

$$S_p(N) = \frac{1}{1 + \alpha_p} \left\{ N^2 + N^p + \sum_{j=1}^{p-1} C_j(N) S_j(N-1) - \sum_{j=1}^{p-1} \alpha_j S_j(N) \right\}. \quad (2)$$

The  $\alpha_j$  in (2) are the coefficients from the  $(p-1)$ -st polynomial solution  $S_{p-1}(N)$ , while the terms  $C_j(N)$  are defined in terms of binomial coefficients<sup>1</sup> as:

$$C_j(N) = \left[ \binom{p-1}{j} N - \binom{p-1}{j-1} \right].$$

2. The closed form solution of  $S_p(N)$  is a polynomial in  $N$  of degree  $p+1$  with rational coefficients and a constant term equal to zero.

While the expression for  $S_p(N)$  in (2) is indeed accurate, it requires repeated iteration to obtain the closed form expression for any particular  $p$ . Such an approach without the assistance of a computer algebra system such as Maxima<sup>2</sup> would be prohibitively time-consuming and prone to error.

We are hence left with the following question:

Can we find a closed-form solution for the coefficients of the general solution polynomial that can be obtained *directly*, i.e. that does not require iteration?

This paper uses linear algebra and matrices to achieve precisely this goal.

## The Direct Approach

### Polynomial with Undetermined Coefficients

Since we have established in (2) that  $S_p(N)$  is a polynomial of order  $p+1$  with no constant term, we may write down the closed form solution of  $S_p(N)$  in the form:

$$S_p(N) := \sum_{j=1}^{p+1} a_j N^j \quad (3)$$

where  $a_j$ , ( $j = 1, 2, \dots, p+1$ ), are coefficients that have yet to be determined. As a first step towards determining these coefficients, observe that every finite summation can be written as a first order recurrence in  $N$  by peeling off the last summand:

$$S_p(N+1) = S_p(N) + (N+1)^p. \quad (4)$$

Substituting (3) into (4) gives:

$$\sum_{j=1}^{p+1} a_j (N+1)^j = \sum_{j=1}^{p+1} a_j N^j + (N+1)^p. \quad (5)$$

<sup>1</sup>The notation  $\binom{n}{k}$  used above denotes binomial coefficients, often expressed verbally as “ $n$  choose  $k$ ”. Other representations of these coefficients include  $C(n, k)$  and  $C_k^n$ .

<sup>2</sup>Maxima code for automatically computing solutions to (1) using solution formula (2) is given in Part 2 of this paper [Ebr10b].

## Summation Manipulations

What follows is a sequence of summation manipulations<sup>3</sup> and simplifications of (5) aimed at obtaining an expression from which a closed-form solution for the coefficients  $a_j$  becomes transparent:

- I Expand the binomial powers using the binomial formula on both the left-hand side (LHS) and right-hand side (RHS) of (5):

$$\sum_{j=1}^{p+1} a_j \sum_{i=0}^j \binom{j}{i} N^i = \sum_{j=1}^{p+1} a_j N^j + \sum_{j=0}^p \binom{p}{j} N^j. \quad (6)$$

- II Modify the LHS and RHS of (6) as follows, exploiting the fact that  $\binom{n}{k} = 0$  when  $k > n$ :

**LHS:** Extend upper range of inner sum to match upper range of outer sum.

**RHS:** Peel off 0-index term and add a vacuous  $p + 1$ -index term to last sum.

**Result:**

$$\sum_{j=1}^{p+1} a_j \sum_{i=0}^{p+1} \binom{j}{i} N^i = \sum_{j=1}^{p+1} a_j N^j + 1 + \sum_{j=1}^{p+1} \binom{p}{j} N^j$$

- III **LHS:** Interchange order of summations.

**RHS:** Combine terms.

**Result:**

$$\sum_{i=0}^{p+1} N^i \sum_{j=1}^{p+1} a_j \binom{j}{i} = 1 + \sum_{j=1}^{p+1} \left[ a_j + \binom{p}{j} N^j \right]$$

- IV **LHS:** Peel off 0-index term from the (now outside) summation on  $i$ .

**RHS:** Change dummy index variable.

**Result:**

$$\sum_{j=1}^{p+1} a_j + \sum_{i=1}^{p+1} N^i \sum_{j=1}^{p+1} a_j \binom{j}{i} = 1 + \sum_{i=1}^{p+1} \left[ a_i + \binom{p}{i} N^i \right]$$

- V Bring everything over to the left-hand side to obtain a homogeneous linear combination of the monomials  $\{N^i\}_{i=0}^{p+1}$ :

$$-1 + \sum_{j=1}^{p+1} a_j + \sum_{i=1}^{p+1} \left[ \left( \sum_{j=1}^{p+1} a_j \binom{j}{i} \right) - a_i - \binom{p}{i} N^i \right] N^i = 0 \quad (7)$$

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<sup>3</sup>Such manipulations are taught in material such as [GKP] and [Knu].

## Linear Independence Reveals Equations for Unknown Coefficients

Since the monomials  $\{N^i\}_{i=0}^{p+1}$  are a basis for the vector space  $\mathbb{P}_{N+1}$  of polynomials of degree less than or equal to  $N + 1$ , *any linear combination of the  $\{N^i\}_{i=0}^{p+1}$  will equal zero if and only if each coefficient equals zero.*

Therefore, setting the coefficients from (7) equal to zero gives us a system of  $p + 2$  simultaneous linear equations in the  $p + 1$  unknown coefficients  $a_j$ , one equation for each of the coefficients of  $N^i$ :

$$(i = 0) \quad \sum_{j=1}^{p+1} a_j = 1 \quad (8)$$

$$(i = 1, \dots, p) \quad \sum_{j=1, j \neq i}^{p+1} \binom{j}{i} a_j = \binom{p}{i} \quad (9)$$

$$(i = p + 1) \quad \sum_{j=1, j \neq p+1}^{p+1} \binom{j}{i} a_j = \binom{p}{i} \quad (*)$$

Observe that (\*) reduces to  $0 = 0$ , since  $j < p + 1$  for all  $j$ , and  $\binom{n}{k} = 0$  when  $n < k$ . Thus we are left with a square system of  $p + 1$  equations ( $i = 0, \dots, p$ ) in  $p + 1$  unknowns  $a_j$ , ( $j = 1, \dots, p + 1$ ).

By observing that  $\binom{j}{i} = 1$  for  $i = 0$ , we can combine (8) and (9) into the single set of linear equations:

$$(i = 0, \dots, p) \quad \sum_{j=1, j \neq i}^{p+1} \binom{j}{i} a_j = \binom{p}{i} \quad (10)$$

Taking into account the inequality  $j \neq i$ , and noting once again that  $\binom{n}{k} = 0$  when  $n < k$ , (10) simplifies to a triangular linear system:

$$(i = 0, \dots, p) \quad \sum_{j=i+1}^{p+1} \binom{j}{i} a_j = \binom{p}{i} \quad (11)$$

and hence all  $p + 1$  equations in the system are linearly independent.

## From Linear System to Matrix Equation

The triangular linear system (11) may be readily solved by combining the  $p + 1$  summations into a single matrix equation and using any one of a number of matrix solver packages. But this requires that the system (11) be re-indexed<sup>4</sup> so that the start value of index  $i$  matches that of index  $j$ . Following the usual matrix convention, referred to as 1-indexing, we choose the start indices to be  $i = j = 1$ .

The 1-indexed system is:

$$(i = 1, \dots, p + 1) \quad \sum_{j=i}^{p+1} \binom{j}{i-1} a_j = \binom{p}{i-1} \quad (12)$$

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<sup>4</sup>Index manipulation is covered in [GKP] and [Knu].

Maxima code for solving the 0-indexed<sup>5</sup> triangular linear system (13) is given in Appendix B. Octave/Matlab code for solving the 1-indexed triangular linear system (12) is given in Appendix C.

The matrix equation equivalent to (12) is:  $\mathbf{M} \cdot \mathbf{a} = \mathbf{b}$ , where

$$\mathbf{M} := \begin{bmatrix} \binom{1}{0} & \binom{2}{0} & \cdots & \binom{p+1}{0} \\ & \binom{2}{1} & \cdots & \binom{p+1}{1} \\ & & \cdots & \binom{p+1}{1} \\ & & \cdots & \binom{p+1}{p} \end{bmatrix}, \quad \mathbf{a} := \begin{bmatrix} a_1 \\ a_2 \\ \cdots \\ a_{p+1} \end{bmatrix}, \quad \text{and} \quad \mathbf{b} := \begin{bmatrix} \binom{p}{0} \\ \binom{p}{1} \\ \cdots \\ \binom{p}{p} \end{bmatrix}. \quad (14)$$

Since  $\mathbf{M}$  is an upper triangular matrix with all non-zero upper triangular entries, we know we can readily solve this for any given values of  $p$  and  $N$  using a matrix solver system such as Maxima/Mathematica, Octave/Matlab, or Maple, among others. Source code for solutions in Maxima and Octave/Matlab are given in Appendix B and Appendix C, respectively.

### Computing the General, Closed-Form Solution

We know from (2) that the general, closed-form solution to (1) is a  $p+1$ -degree polynomial (3) in  $N$  with rational coefficients  $a_j$  given by solving the triangular linear system (12), or equivalently (14).

The highest few coefficients  $a_j$  can be readily computed by hand. For all  $p$ , we have:

1.  $a_{p+1} = \frac{1}{p+1}$
2.  $a_p = \frac{1}{2} \binom{p}{0} = \frac{1}{2}$
3.  $a_{p-1} = \frac{1}{12} \binom{p}{1} = \frac{p}{12}$
4.  $a_{p-3} = -\frac{1}{120} \binom{p}{3}$
5.  $a_{p-5} = \frac{1}{252} \binom{p}{5}$
6.  $a_{p-7} = -\frac{1}{240} \binom{p}{7}$
7.  $a_{p-9} = \frac{1}{132} \binom{p}{9}$ .
8. In particular, we claim (without proof) that

$$a_{p-i} = c_i \binom{p}{i}, \quad (15)$$

where the  $c_i$  are rational coefficients independent of  $p$ , and  $c_i = 0$  for all even positive  $i$ . Hence we have:

$$a_{p-2} = a_{p-4} = \dots = 0.$$

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<sup>5</sup>0-indexing is more convenient for various computational software packages and programming languages, including Maxima. The 0-indexed system is:

$$(i = 0, \dots, p) \quad \sum_{j=i}^p \binom{j+1}{i} a_j = \binom{p}{i} \quad (13)$$

Substituting the above coefficients into the general closed-form polynomial solution gives us:

$$S_p(N) = \frac{1}{p+1}N^{p+1} + \frac{1}{2}N^p + \frac{p}{12}N^{p-1} - \frac{1}{120}\binom{p}{3}N^{p-3} + \frac{1}{252}\binom{p}{5}N^{p-5} \\ - \frac{1}{240}\binom{p}{7}N^{p-7} + \frac{1}{132}\binom{p}{9}N^{p-9} - \dots + pc_{p-1}N. \quad (16)$$

A listing of the solution formulas for the first ten  $p$  values is given in Appendix A.

## Conclusion

Reviewing the path we have taken in this 3-part paper: Part 1 [Ebr10a] motivated the problem and illustrated a recurrence relation method for obtaining the solution for small  $p$ . Part 2 [Ebr10b] generalized this method and found a  $p$ -th order recurrence relation for the solution. Part 2 also illustrated, by an induction argument, that the closed form solutions for all  $p$  are polynomials of degree  $p + 1$  with rational coefficients and no constant term.

This paper used the closed-form polynomial expression motivated in Part 2 to obtain a direct solution. By writing out the polynomial form with undetermined coefficients, the recurrence relation (4) was manipulated into a linear combination of polynomials (7). Using a linear independence argument, (7) was reduced to a triangular linear system (12) and associated matrix equation (14). The general closed-form solution is (16).

Solving the finite-summation-of-integer-powers problem  $\sum_{k=1}^N k^p$ , for arbitrary positive integers  $p$  and  $N$  has provided a natural setting to use a variety of techniques from discrete mathematics and linear algebra, and poses additional interesting questions (characterization of the denominators in the coefficients of (16) (the  $c_i$  in (15)) and divisibility properties of  $S_p$  for various  $p$ ). In particular, we have obtained a direct method for solving the finite-summation-of-integer-powers problem that invokes a matrix solution and does not require iteration.

## Appendix A: Listing of Solutions to (1)

A listing of the solution formulas for the first ten  $p$  values is as follows:

$$\begin{aligned}
 S_1(N) &= \sum_{k=1}^N k = \frac{N^2}{2} + \frac{N}{2} \\
 S_2(N) &= \sum_{k=1}^N k^2 = \frac{N^3}{3} + \frac{N^2}{2} + \frac{N}{6} \\
 S_3(N) &= \sum_{k=1}^N k^3 = \frac{N^4}{4} + \frac{N^3}{2} + \frac{N^2}{4} \\
 S_4(N) &= \sum_{k=1}^N k^4 = \frac{N^5}{5} + \frac{N^4}{2} + \frac{N^3}{3} - \frac{N}{30} \\
 S_5(N) &= \sum_{k=1}^N k^5 = \frac{N^6}{6} + \frac{N^5}{2} + \frac{5N^4}{12} - \frac{N^2}{12} \\
 S_6(N) &= \sum_{k=1}^N k^6 = \frac{N^7}{7} + \frac{N^6}{2} + \frac{N^5}{2} - \frac{N^3}{6} + \frac{N}{42} \\
 S_7(N) &= \sum_{k=1}^N k^7 = \frac{N^8}{8} + \frac{N^7}{2} + \frac{7N^6}{12} - \frac{7N^4}{24} + \frac{N^2}{12} \\
 S_8(N) &= \sum_{k=1}^N k^8 = \frac{N^9}{9} + \frac{N^8}{2} + \frac{2N^7}{3} - \frac{7N^5}{15} + \frac{2N^3}{9} - \frac{N}{30} \\
 S_9(N) &= \sum_{k=1}^N k^9 = \frac{N^{10}}{10} + \frac{N^9}{2} + \frac{3N^8}{4} - \frac{7N^6}{10} + \frac{N^4}{2} - \frac{3N^2}{20} \\
 S_{10}(N) &= \sum_{k=1}^N k^{10} = \frac{N^{11}}{11} + \frac{N^{10}}{2} + \frac{5N^9}{6} - N^7 + N^5 - \frac{N^3}{2} + \frac{5N}{66}
 \end{aligned}$$

## Appendix B: Maxima Source Code for Solving (14)

For particular  $p$ , the coefficients  $a_j$  for the closed form solutions  $S_p(N)$  are immediate. Maxima code, is given below, for determining this solution using the 0-indexed triangular linear system (13):

```

solution(p):= block([a, eq], /* give subroutine variables local scope */
  v : makelist(a[i], i, 0, p), /* create list of unknowns (0-indexed) */
                                /* create list of equations (0-indexed) */
  eq : makelist(sum(binom(j+1,i)*a[j],j,i,p) = binom(p,i), i, 0, p),
  linsolve(eq, v)
)$

```

A more elaborate function that computes the series form and factored forms of the solution is given below. This was used to generate the first ten solution listings.

```

SpN_mat(p):= block([a, eq], /* give subroutine variables local scope */
  v : makelist(a[i], i, 0, p), /* create list of unknowns (0-indexed) */
                                /* create list of equations (0-indexed) */

```

```

    eq : makelist(sum(binom(j+1,i)*a[j],j,i,p) = binom(p,i), i, 0, p),
/* find coefficients of solution polynomial by solving linear system */
    sol : linsolve(eq, v),
/* create polynomial: inner product of {N^i} with coefficients */
    pol : makelist(N^(i+1),i,0,p), /* monomials {N^i} */
/* closed form formula */
    cff : rhs(sol.pol), /* inner product of {N^i} with coefficients */
    cff /* return closed form formula in series */
)$

```

## Appendix C: Octave/Matlab Source Code for Solving (14)

For particular  $p$ , the coefficients  $a_j$  for the closed form solutions  $S_p(N)$  are immediate. Octave/Matlab code, is given below, for determining this solution using the 1-indexed triangular linear system (12):

```

% Function listing coefficients from a_p+1 to a_1 as fractions
function c = sumkp_matrix(p)
    M=zeros(p+1); % initialize matrix M
    for i=1:p+1
        for j=i:p+1 % set upper triangular elements
            M(i,j)=nchoosek(j,i-1); % Equation (12)
        end;
    end;

    b=zeros(p+1,1); % initialize column matrix b
    for i=1:p+1
        b(i)=nchoosek(p,i-1); % Equation (12)
    end

    c=inv(M)*b; % Solve to obtain coefficients

    c=flipr(c'); % List from highest index to lowest
    sml=abs(c)<1e-8; % Numerical correction: set tiny coefficients to exactly zero
    c(find(sml==1))=0;

    disp("Solution Coefficients (High to Low order):")
    c=rats(c); % return coefficients as fractions
end;

```

## References

- [Ebr10a] Assad Ebrahim. Finite summation of integer powers  $x^p$ , part 1. January 2010.
- [Ebr10b] Assad Ebrahim. Finite summation of integer powers  $x^p$ , part 2. February 2010.
- [GKP] Ron Graham, Donald Knuth, and Oren Patashnik. *Concrete Mathematics: A Foundation for Computer Science*. Addison Wesley.
- [Knu] Donald Knuth. *The Art of Computer Programming (3 Volumes)*.